

Logic programming III

Henrik Boström
Stockholm University

- Arithmetics
- Meta-logical predicates
- Cut
- Extra-logical predicates
- Predicates to find all solutions

Arithmetics

- | | |
|----------------------|---|
| $X \text{ is } Y$ | - The arithmetic expression Y is evaluated and unified with X . |
| $X =: = Y$ | - The values of X and Y are the same. |
| $X = \backslash = Y$ | - The values of X and Y are different. |
| $X < Y$ | - The value of X is less than the value of Y . |
| $X > Y$ | - The value of X är is greater than the value of Y . |
| $X = < Y$ | - The value of X is less than or equal to the value of Y . |
| $X > = Y$ | - The value of X is greater than or equal to the value of Y . |

Example

```
sum([],0).  
sum([X|L],Sum):-  
    sum(L,RestSum),  
    Sum is RestSum+X.
```

```
sum(L,Sum):- sum(L,0,Sum).
```

```
sum([],Sum,Sum).  
sum([X|L],PreviousSum,Sum):-  
    SumSoFar is PreviousSum+X,  
    sum(L,SumSoFar,Sum).
```

Meta-logical predicates

Term1 == Term2	- Term1 and Term2 are identical
Term1 \== Term2	- Term1 and Term2 are not identical
Term1 = Term2	- Term1 and Term2 are unified
var(X)	- X is an (uninstantiated) variable
nonvar(X)	- X is not an (uninstantiated) variable
ground(X)	- X does not contain a variable
atom(X)	- X is a constant and not a number
number(X)	- X is a number
atomic(X)	- X is a constant or a number
compound(X)	- X is a term with arity > 0

Metalogical predicates

$\text{functor}(T, N, A)$ - The term T has the term name N and the arity A .

?- $\text{functor}(p(a, X), N, A)$.

$N = p \quad A = 2$

?- $\text{functor}(T, f, 3)$.

$T = f(X, Y, Z)$

$\text{arg}(\text{No}, T, \text{Arg})$ - Argument no. No in the term T is unified with Arg .

?- $\text{arg}(3, p(0, s(0), s(s(0))), X)$.

$X = s(s(0))$

?- $\text{arg}(2, p(a, b), a)$.

no

Meta-logical predicates

$T =.. L$ - L is a list where the first element is the term name of T and the rest of the list consists of the arguments of T .

?- $p(a, s(X), Y) =.. L$.

$L = [p, a, s(X), Y]$

?- $T =.. [a, b, c, d]$.

$T = a(b, c, d)$

Unification with occur's check

```
unify(X,Y):- var(X), var(Y), X=Y.  
unify(X,Y):- atomic(X), atomic(Y), X=Y.  
unify(X,Y):- var(X), nonvar(Y), not_in(X,Y), X=Y.  
unify(X,Y):- nonvar(X), var(Y), not_in(Y,X), X=Y.  
unify(X,Y):- compound(X),compound(Y), term_unify(X,Y).  
  
not_in(X,Y):- var(Y), X \== Y.  
not_in(X,Y):- nonvar(Y), functor(Y,_,N), not_in(X,N,Y).  
  
not_in(_,0,_).  
not_in(X,N1,Y):- N1 > 0, arg(N1,Y,Arg), not_in(X,Arg),  
N2 is N1-1, not_in(X,N2,Y).
```

Unification with occur's check (cont.)

```
term_unify(X,Y):-  
    functor(X,F,N),  
    functor(Y,F,N),  
    unify_args(N,X,Y).  
  
unify_args(0,_,_).  
unify_args(N1,X,Y):-  
    N1 > 0,  
    arg(N1,X,ArgX), arg(N1,Y,ArgY),  
    unify(ArgX,ArgY), N2 is N1-1,  
    unify_args(N2,X,Y).
```

Cut (!)

Cut is a predicate that is used to reduce the number of (unnecessary) goal reductions. Assume we have a unary goal G and a clause $G' :- G_1, \dots, G_{i-1}, !, G_{i+1}, \dots, G_n$, where G and G' have an mgu θ .

If cut has been reduced (after $(G_1, \dots, G_{i-1})\theta$) the following holds:

- No other clause will be used for reducing G .
- Alternative reductions for $(G_1, \dots, G_{i-1})\theta$ will not be tested.
- The remaining goals $(G_{i+1}, \dots, G_n)\theta$ are not affected by the cut.

Example (green cut)

```
merge([X|Xs],[Y|Ys],[X|Zs]):-
```

```
  X =< Y, !,
```

```
  merge(Xs,[Y|Ys],Zs).
```

```
merge([X|Xs],[Y|Ys],[Y|Zs]):-
```

```
  X > Y, !,
```

```
  merge([X|Xs],Ys,Zs).
```

```
merge([],Ys,Ys):-!.
```

```
merge(Xs,[],Xs).
```

Exempel (red cut)

```
delete(_, [], []).  
delete(X, [X|L1], L2): -  
    delete(X, L1, L2).  
delete(X, [Y|L1], [Y|L2]): -  
    X \== Y,  
    delete(X, L1, L2).
```

```
delete(_, [], []).  
delete(X, [X|L1], L2): - !, delete(X, L1, L2).  
delete(X, [Y|L1], [Y|L2]): - delete(X, L1, L2).
```

Example

```
min(X, Y, X): - X =< Y.  
min(X, Y, Y): - X > Y.
```

```
min(X, Y, X): - X =< Y, !.  
min(X, Y, Y).
```

```
min(X, Y, Z): - X =< Y, !, X=Z.  
min(X, Y, Y).
```

Example

```
if_then_else(P,Q,_):- P, Q.  
if_then_else(P,_,R):- \+ P, R.
```

```
if_then_else(P,Q,_):- P, !, Q.  
if_then_else(_,_,R):- R.
```

```
not(G):- G, !, fail.  
not(_).
```

Extralogical predicates

read(X)	- read term from current stream
write(X)	- write term to current stream
see(File)	- open the file File for reading
seen	- close the current file
tell(File)	- open the file File for writing
told	- close the current file

```
?- write('Give answer: '), read(Answer).
```

```
Give answer: yes.
```

```
Answer = yes
```

Example

```
process_data(Input,Output): -  
    see(Input),  
    tell(Output),  
    repeat,  
    read(X),  
    do_something(X),  
    X == end_of_file,  
    told,  
    seen.
```

```
do_something(end_of_file): -!.
```

```
do_something(X): -  
    transform(X,Y),  
    write(Y), write('.'),nl.
```

Extra-logical predicates

- assert(C) - The clause C is added to memory
- asserta(C) - C is added first
- assertz(C) - C is added last
- retract(H)* - The first clause with a head unifying with H is removed.
- retractall(H) - All clauses having a head that unifies with H are removed
- clause(H,B)* - Unify H with the head and B with the body of a rule.

*More than one answer can be obtained (upon back-tracking).

Example

$p(a). p(b). p(c).$
 $q(1). q(2). q(3).$

?- $p(X), q(Y), \text{assert}(r(X,Y)), \text{fail}.$

$r(a,1). r(a,2). r(a,3).$
 $r(b,1). r(b,2). r(b,3).$
 $r(c,1). r(c,2). r(c,3).$

Findall

$\text{findall}(E,G,L)$

- for every reduction of G, add element E to L

$\text{parent}(a,b). \text{parent}(e,b). \text{parent}(b,c). \text{parent}(b,d).$

?- $\text{findall}(P,\text{parent}(P,C),L).$

$L = [a,e,b,b]$

?- $\text{findall}((GP,GC), (\text{parent}(GP,P), \text{parent}(P,GC)),L).$

$L = [(a,c),(a,d),(e,c),(e,d)]$

Bagof

bagof(E,G,L)

- For each reduction of G, where non-quantified (i.e., they are not in E or specified with \wedge) are assumed to refer to the same instance, add E to L.

parent(a,b). parent(e,b). parent(b,c). parent(b,d).

?- bagof(P,parent(P,C),L).

C = b L = [a,e];

C = c L = [b];

C = d L = [b]

?- bagof(P,C \wedge parent(P,C),L).

L = [a,e,b,b]

Setof

setof(E,G,L)

- as bagof(E,G,L), but where L is sorted and without duplicates.

parent(a,b). parent(e,b). parent(b,c). parent(b,d).

?- setof(P,parent(P,C),L).

C = b L = [a,e];

C = c L = [b];

C = d L = [b]

?- setof(P,C \wedge parent(P,C),L).

L = [a,b,e]